

# Uncloneable states are necessary as proofs and advice

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- Secure software leasing: Ananth and La Placa (2021), Broadbent, Jeffery, Lord, Podder and Sundaram (2021), ...

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2. Same as 1. with **advice**.
3. There is a relational problem which can only be **exactly** solved by a quantum poly-time algorithm with an uncloneable quantum state as quantum **advice**.

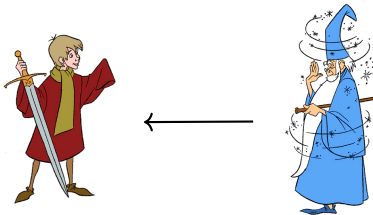
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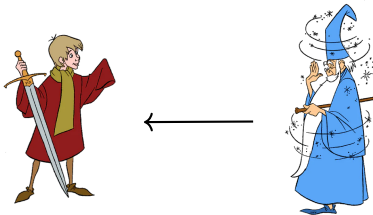
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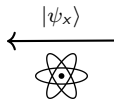
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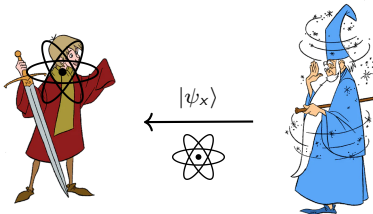
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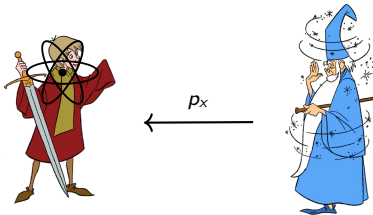
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- ▶ Classical proofs are always cloneable
- ▶ We use the **AK07** quantum oracle

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- ▶ **FBQP/qpoly**: like BQP/qpoly for polynomially bounded relations  $R \subseteq \{0, 1\}^* \times \{0, 1\}^*$ , ...
- ▶ In fact the **ABK23** construction separates **FEQP/qpoly** and FBQP/poly
- ▶ We use the **ABK23** construction

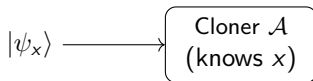
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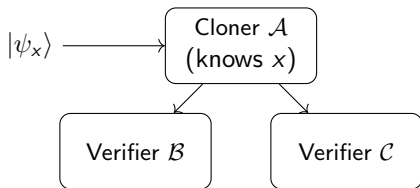
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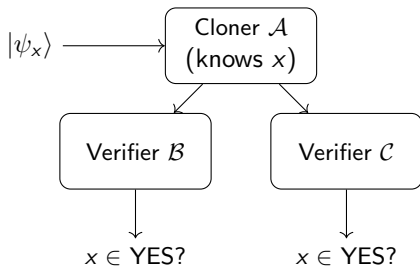
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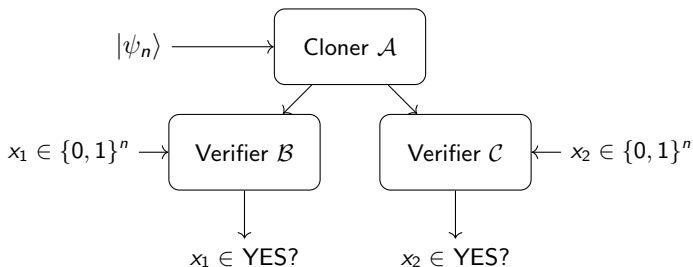
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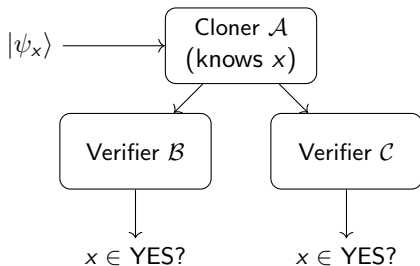
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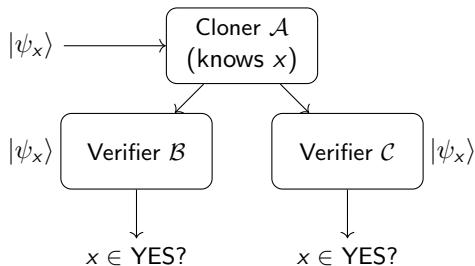


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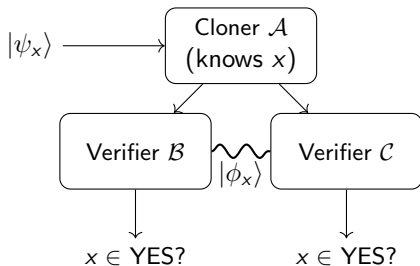
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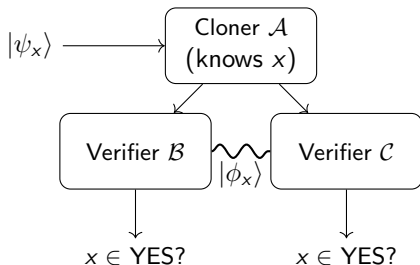
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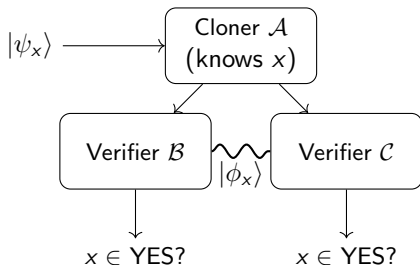
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- Often average-case uncloneability is considered

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- ▶ Uncloneability is **worst-case** instead of average case

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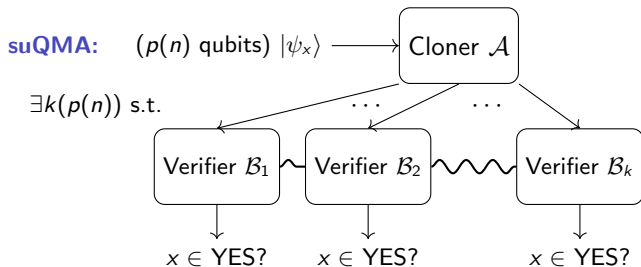
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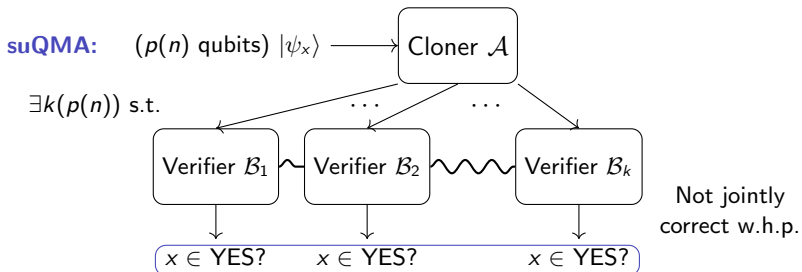
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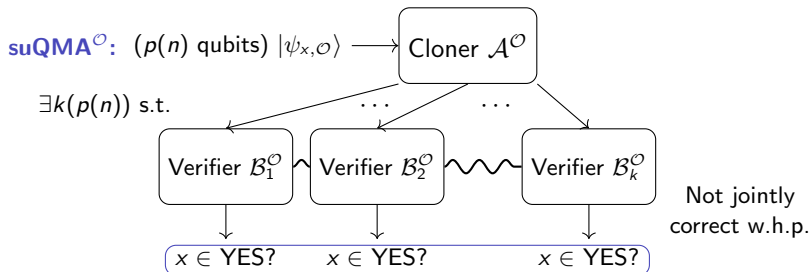




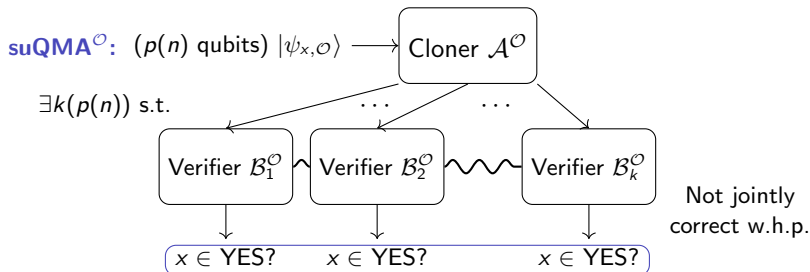
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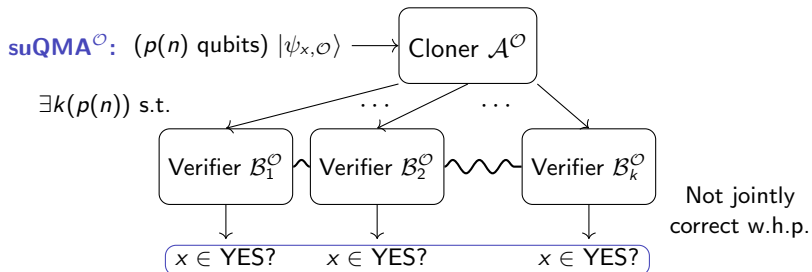


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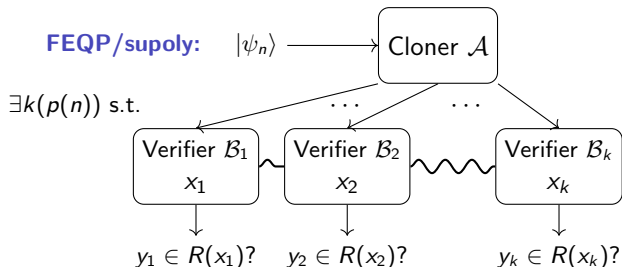
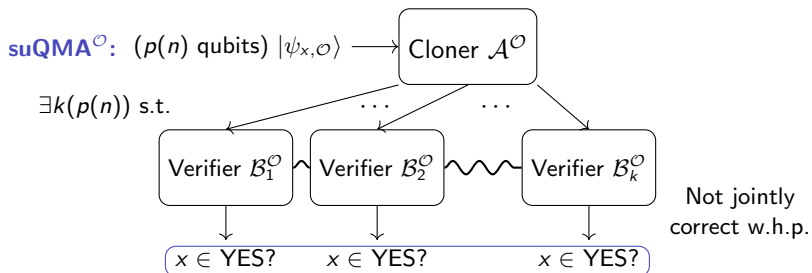
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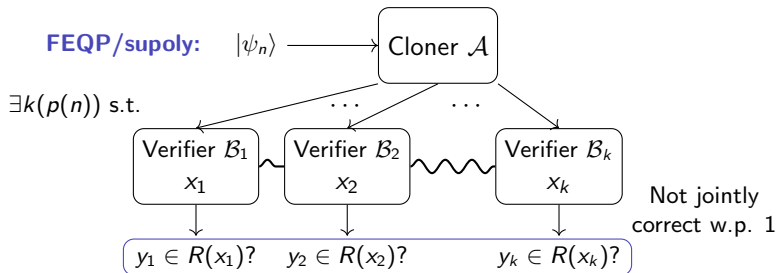
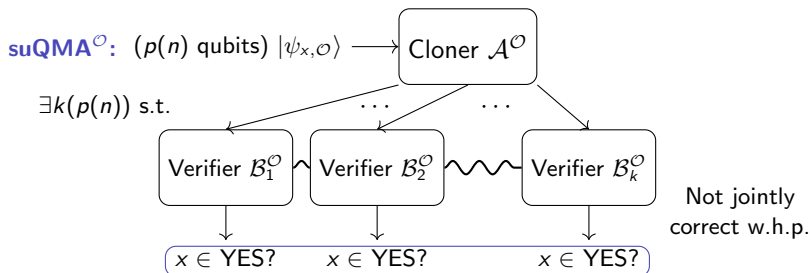


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**Canonical advice:**  $\frac{1}{2^{n/2}} \sum_{x \in \{0, 1\}^n} (-1)^{f(x)} |x\rangle$

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**ABK23** problem (Hidden Matching): Family of relations  $\{R_f\}_f \subseteq \{0, 1\}^n \times \{0, 1\}^{n+1}$  indexed by  $f : \{0, 1\}^n \rightarrow \{0, 1\}$

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**Canonical advice:**  $\frac{1}{2^{n/2}} \sum_{x \in \{0, 1\}^n} (-1)^{f(x)} |x\rangle$

Good news: Binary phase states (for random  $f$ ) are uncloneable!

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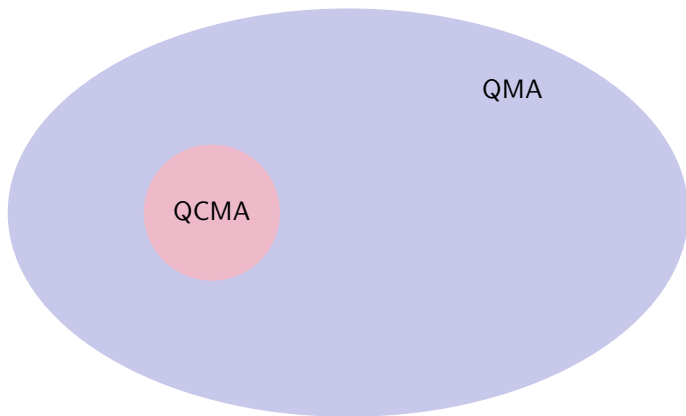
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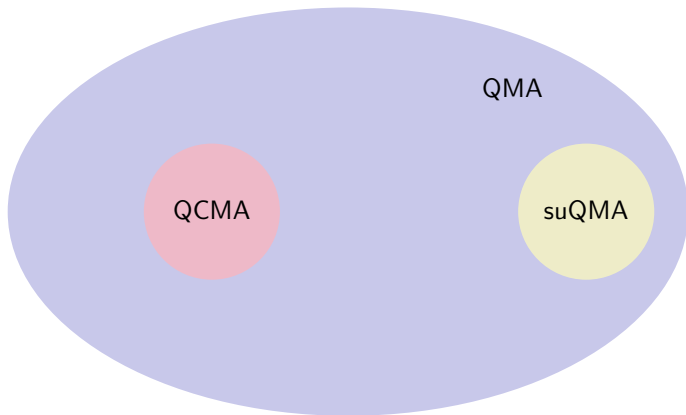
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- ▶ For FEQP/supoly, we show:
  - Proofs  $|\psi_f\rangle$  that works jointly for  $\mathcal{B}_1, \dots, \mathcal{B}_k$  must have inner product  $|\langle\psi_{f'}|\psi_f\rangle| = 2^{-\Omega(k)}$ ;
  - In order to achieve the inner products, the states need  $\Omega(k)$  qubits.

# Other consequences

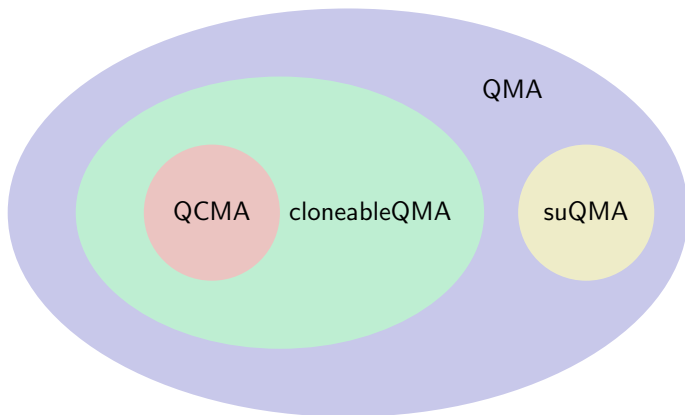
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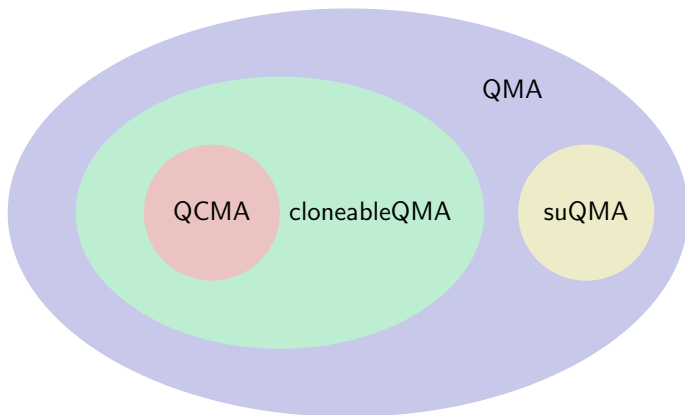
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From definition,  $\text{suQMA} \subseteq \text{QMA} \setminus \text{cloneableQMA}$   
 $\Rightarrow$  w.r.t. a quantum oracle,  $\text{QMA} \neq \text{cloneableQMA}$

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